

# DATABASE SYSTEMS

DESIGN IMPLEMENTATION AND MANAGEMENT

INTERNATIONAL EDITION



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## Chapter 7 Normalisation

## In this chapter, you will learn:

- What normalization is and what role it plays in the database design process
- About the normal forms 1NF, 2NF, 3NF, BCNF, and 4NF
- How normal forms can be transformed from lower normal forms to higher normal forms
- That normalization and ER modeling are used concurrently to produce a good database design
- That some situations require denormalization to generate information efficiently



# Database Tables and Normalization

- Normalization
  - Process for evaluating and correcting table structures to minimize data redundancies
    - Reduces data anomalies
  - Works through a series of stages called normal forms:
    - First normal form (1NF)
    - Second normal form (2NF)
    - Third normal form (3NF)



# Database Tables and Normalization (continued)

- Normalization (continued)
  - 2NF is better than 1NF; 3NF is better than 2NF
  - For most business database design purposes, 3NF is as high as we need to go in normalization process
  - Highest level of normalization is not always most desirable



## The Need for Normalization

- Example: Company that manages building projects
  - Charges its clients by billing hours spent on each contract
  - Hourly billing rate is dependent on employee's position
  - Periodically, report is generated that contains information displayed in Table 7.1



## The Need for Normalization

**TABLE 7.1** A sample report layout

Proj. Num.	Project Name	Employee Number	Employee Name	Job Class	Chg/ Hour	Hours Billed	Total Charge
15	Evergreen	103	June E. Arbough	Elec. Engineer	€67.55	23.8	€1,607.69
		101	John G. News	Database Designer	€82.95	19.4	€1,609.23
		105	Alice K. Johnson*	Database Designer	€82.95	35.7	€2,961.32
		106	William Smithfield	Programmer	€26.66	12.6	€335.92
		102	David H. Senior	Systems Analyst	€76.43	23.8	€1,819.03
				<b>Subtotal</b>			<b>€8,333.19</b>
18	Amber Wave	114	Annelise Jones	Applications Designer	€38.00	25.6	€972.80
		118	James J. Frommer	General Support	€14.50	45.3	€656.85
		104	Anne K. Ramoras*	Systems Analyst	€76.43	32.4	€2,476.33
		112	Darlene M. Smithson	DSS Analyst	€36.30	45.0	€1,633.50
				<b>Subtotal</b>			<b>€5,739.48</b>
22	Rolling Tide	105	Alice K. Johnson	Database Designer	€82.95	65.7	€5,449.82
		104	Anne K. Ramoras	Systems Analyst	€76.43	48.4	€3,699.21
		113	Delbert K. Joenbrood*	Applications Designer	€38.00	23.6	€896.80
		111	Geoff B. Wabash	Clerical Support	€21.23	22.0	€467.06
		106	William Smithfield	Programmer	€28.24	12.8	€361.47
				<b>Subtotal</b>			<b>€10,874.36</b>
25	Starflight	107	Maria D. Alonzo	Programmer	€28.24	25.6	€722.94
		115	Travis B. Bawangi	Systems Analyst	€76.43	45.8	€3,500.49
		101	John G. News*	Database Designer	€82.95	56.3	€4,670.09
		114	Annelise Jones	Applications Designer	€38.00	33.1	€1,257.80
		108	Ralph B. Washington	Systems Analyst	€76.43	23.6	€1,803.75
		118	James J. Frommer	General Support	€14.50	30.5	€442.25
		112	Darlene M. Smithson	DSS Analyst	€36.30	41.4	€1,502.82
				<b>Subtotal</b>			<b>€13,900.14</b>
				<b>Total</b>			<b>€38,942.09</b>

Note: \* indicates project leader.



# The Need for Normalization

**FIGURE 7.1** Tabular representation of the report format

Database name: Ch07\_ConstructCo

Table name: RPT\_FORMAT

RPT_FORMAT						
PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
15	Evergreen	103	June E. Arbough	Elect. Engineer	€67.55	23.80
		101	John G. News	Database Designer	€62.95	19.40
		105	Alice K. Johnson *	Database Designer	€62.95	35.70
		106	William Smithfield	Programmer	€26.66	12.60
		102	David H. Senior	Systems Analyst	€76.43	23.80
18	Amber Wave	114	Annelise Jones	Applications Designer	€38.00	24.60
		118	James J. Frommer	General Support	€14.50	45.30
		104	Anne K. Ramoras *	Systems Analyst	€76.43	32.40
		112	Darlene M. Smithson	DSS Analyst	€36.30	44.00
22	Rolling Tide	105	Alice K. Johnson	Database Designer	€62.95	64.70
		104	Anne K. Ramoras	Systems Analyst	€76.43	46.40
		113	Delbert K. Joenbrood *	Applications Designer	€38.10	23.60
		111	Geoff B. Wabash	Clerical Support	€21.23	22.00
		106	William Smithfield	Programmer	€26.24	12.80
25	Starflight	107	Maria D. Alonzo	Programmer	€26.24	24.60
		115	Travis B. Bawangi	Systems Analyst	€76.43	45.80
		101	John G. News *	Database Designer	€62.95	56.30
		114	Annelise Jones	Applications Designer	€38.00	33.10
		108	Ralph B. Washington	Systems Analyst	€76.43	23.60
		118	James J. Frommer	General Support	€14.50	30.50
		112	Darlene M. Smithson	DSS Analyst	€36.30	41.40

## The Need for Normalization (continued)

- Structure of data set in Figure 7.1 does not handle data very well
- The table structure appears to work; report generated with ease
- Unfortunately, report may yield different results depending on what data anomaly has occurred





## The Normalization Process

- Each table represents a single subject
- No data item will be unnecessarily stored in more than one table
- All attributes in a table are dependent on the primary key



# The Normalization Process (continued)

**TABLE 7.2** Normal forms

Normal Form	Characteristic	Section
First normal form (1NF)	Table format; no repeating groups and PK identified	7.3.1
Second normal form (2NF)	1NF and no partial dependencies	7.3.2
Third normal form (3NF)	2NF and no transitive dependencies	7.3.3
Boyce-Codd normal form (BCNF)	Every determinant is a candidate key (special case of 3NF)	7.6.1
Fourth normal form (4NF)	3NF and no independent multivalued dependencies	7.6.2



# Conversion to First Normal Form

- Repeating group
  - Derives its name from the fact that a group of multiple entries of same type can exist for any single key attribute occurrence
- Relational table must not contain repeating groups
- Normalizing table structure will reduce data redundancies
- Normalization is three-step procedure



# Conversion to First Normal Form (continued)

- Step 1: Eliminate the Repeating Groups
  - Present data in tabular format, where each cell has single value and there are no repeating groups
  - Eliminate repeating groups, eliminate nulls by making sure that each repeating group attribute contains an appropriate data value



# Conversion to First Normal Form (continued)

**FIGURE 7.2** A table in first normal form

Database name: Ch07\_ConstructCo

Table name: DATA\_ORG\_1NF

DATA_ORG_1NF						
PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
15	Evergreen	103	Jane E. Arbough	Elect. Engineer	€87.55	23.80
15	Evergreen	101	John G. News	Database Designer	€82.95	19.40
15	Evergreen	105	Alice K. Johnson *	Database Designer	€82.95	35.70
15	Evergreen	106	William Smithfield	Programmer	€26.66	12.60
15	Evergreen	102	David H. Senior	Systems Analyst	€76.43	23.80
18	Amber Wave	114	Annelise Jones	Applications Designer	€38.00	24.60
18	Amber Wave	118	James J. Frommer	General Support	€14.50	45.30



# Conversion to First Normal Form (continued)

DATA_ORG_1NF						
PROJ_NUM	PROJ_NAME	EMP_NUM	EMP_NAME	JOB_CLASS	CHG_HOUR	HOURS
18	Amber Wave	104	Anne K. Ramoras *	Systems Analyst	€76.43	32.40
18	Amber Wave	112	Darlene M. Smithson	DSS Analyst	€38.30	44.00
22	Rolling Tide	105	Alice K. Johnson	Database Designer	€82.95	64.70
22	Rolling Tide	104	Anne K. Ramoras	Systems Analyst	€76.43	48.40
22	Rolling Tide	113	Delbert K. Joenbrood *	Applications Designer	€38.00	23.60
22	Rolling Tide	111	Geoff B. Wabash	Clerical Support	€21.23	22.00
22	Rolling Tide	106	William Smithfield	Programmer	€28.24	12.80
25	Starflight	107	Maria D. Alonzo	Programmer	€28.24	24.60
25	Starflight	115	Travis B. Bawangi	Systems Analyst	€76.43	45.60
25	Starflight	101	John G. News *	Database Designer	€82.95	56.30
25	Starflight	114	Annelise Jones	Applications Designer	€38.00	33.10
25	Starflight	108	Ralph B. Washington	Systems Analyst	€76.43	23.60
25	Starflight	118	James J. Frommer	General Support	€14.50	30.50
25	Starflight	112	Darlene M. Smithson	DSS Analyst	€38.30	41.40



# Conversion to First Normal Form (continued)

- Step 2: Identify the Primary Key
  - Primary key must uniquely identify attribute value
  - New key must be composed



# Conversion to First Normal Form (continued)

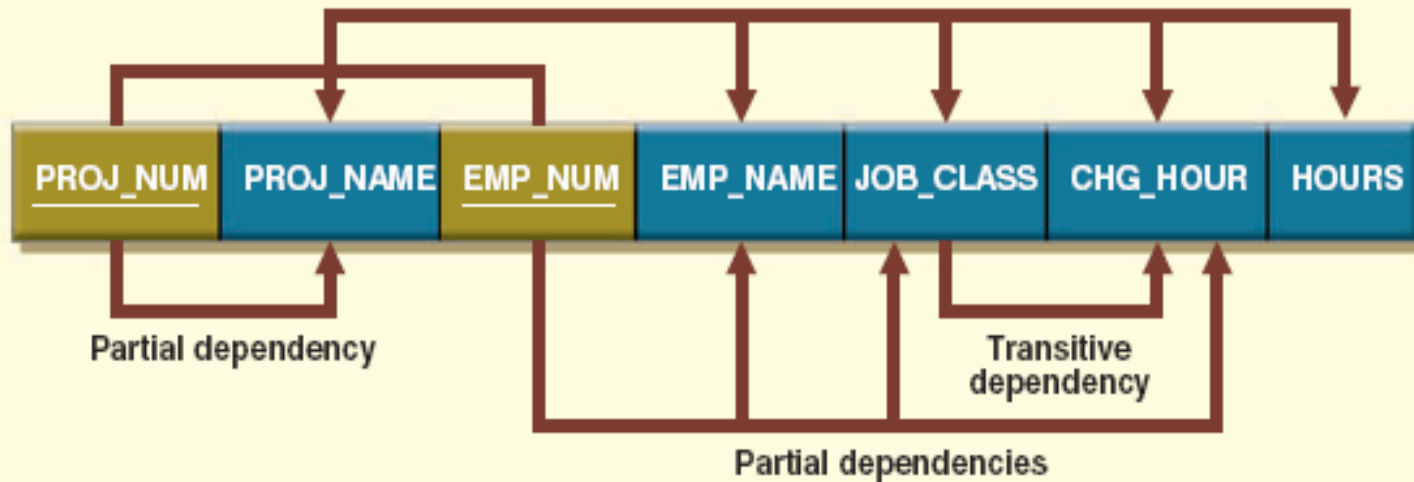
- Step 3: Identify All Dependencies
  - Dependencies can be depicted with help of a diagram
  - Dependency diagram:
    - Depicts all dependencies found within given table structure
    - Helpful in getting bird's-eye view of all relationships among table's attributes
    - Makes it less likely that will overlook an important dependency





# Conversion to First Normal Form (continued)

**FIGURE 7.3** First normal form (1NF) dependency diagram



1NF (PROJ\_NUM, EMP\_NUM, PROJ\_NAME, EMP\_NAME, JOB\_CLASS, CHG\_HOURS, HOURS)

PARTIAL DEPENDENCIES:

(PROJ\_NUM → PROJ\_NAME)

(EMP\_NUM → EMP\_NAME, JOB\_CLASS, CHG\_HOUR)

TRANSITIVE DEPENDENCY

(JOB\_CLASS → CHG\_HOUR)

## Conversion to First Normal Form (continued)

- First normal form describes tabular format in which:
  - All key attributes are defined
  - There are no repeating groups in the table
  - All attributes are dependent on primary key
- All relational tables satisfy 1NF requirements
- Some tables contain partial dependencies
  - Dependencies based on only part of the primary key
  - Sometimes used for performance reasons, but should be used with caution
  - Still subject to data redundancies





## Conversion to Second Normal Form

- Relational database design can be improved by converting the database into second normal form (2NF)
- Two steps





# Conversion to Second Normal Form (continued)

- Step 1: Write Each Key Component on a Separate Line
  - Write each key component on separate line, then write original (composite) key on last line
  - Each component will become key in new table



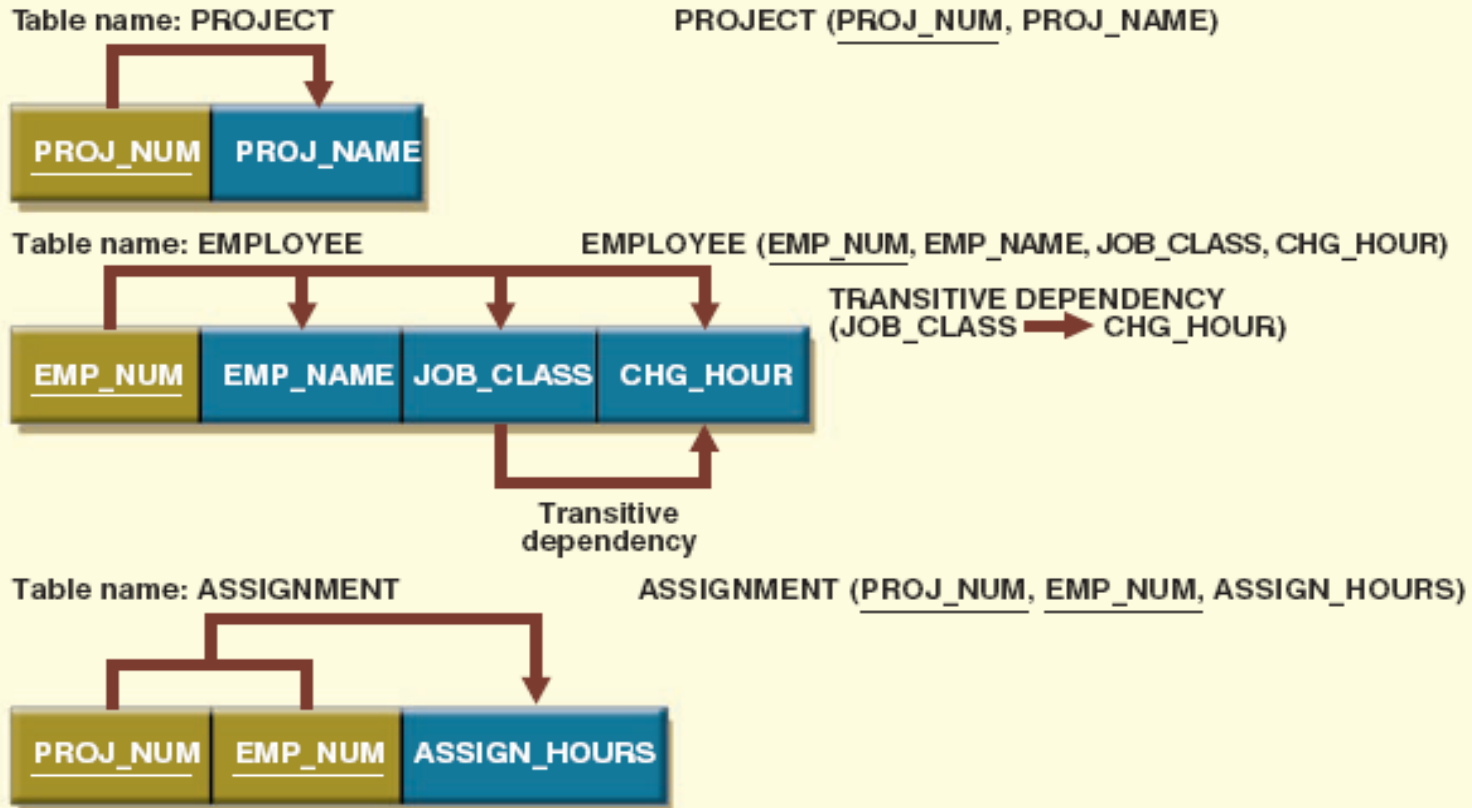
# Conversion to Second Normal Form (continued)

- Step 2: Assign Corresponding Dependent Attributes
  - Determine those attributes that are dependent on other attributes
  - At this point, most anomalies have been eliminated



# Conversion to Second Normal Form (continued)

**FIGURE 7.4** Second normal form (2NF) conversion results



# Conversion to Second Normal Form (continued)

- Table is in second normal form (2NF) when:
  - It is in 1NF and
  - It includes no partial dependencies:
    - No attribute is dependent on only portion of primary key



## Conversion to Third Normal Form

- Data anomalies created are easily eliminated by completing three steps
- Step 1: Identify Each New Determinant
  - For every transitive dependency, write its determinant as PK for new table
    - Determinant
      - Any attribute whose value determines other values within a row





# Conversion to Third Normal Form (continued)

- Step 2: Identify the Dependent Attributes
  - Identify attributes dependent on each determinant identified in Step 1 and identify dependency
  - Name table to reflect its contents and function



## Conversion to Third Normal Form (continued)

- Step 3: Remove the Dependent Attributes from Transitive Dependencies
  - Eliminate all dependent attributes in transitive relationship(s) from each of the tables that have such a transitive relationship
  - Draw new dependency diagram to show all tables defined in Steps 1–3
  - Check new tables as well as tables modified in Step 3 to make sure that each table has determinant and that no table contains inappropriate dependencies



# Conversion to Third Normal Form (continued)

**FIGURE 7.5**

Third normal form (3NF) conversion results



Table name: PROJECT

PROJECT (PROJ\_NUM, PROJ\_NAME)

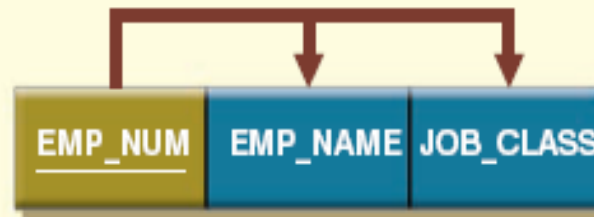


Table name: EMPLOYEE

EMPLOYEE (EMP\_NUM, EMP\_NAME, JOB\_CLASS)



Table name: JOB

JOB (JOB\_CLASS, CHG\_HOUR)

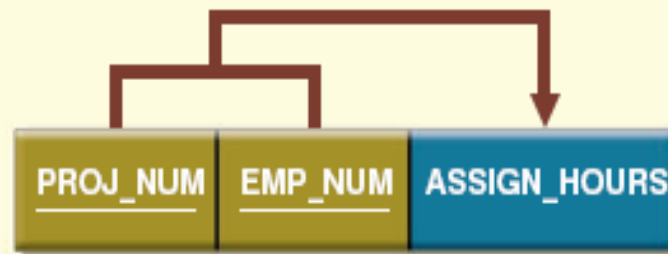


Table name: ASSIGNMENT

ASSIGNMENT (PROJ\_NUM, EMP\_NUM, ASSIGN\_HOURS)



## Conversion to Third Normal Form (continued)

- A table is in third normal form (3NF) when both of the following are true:
  - It is in 2NF
  - It contains no transitive dependencies





## Improving the Design

- Table structures are cleaned up to eliminate troublesome initial partial and transitive dependencies
- Normalization cannot, by itself, be relied on to make good designs
- It is valuable because its use helps eliminate data redundancies



## Improving the Design (continued)

- Issues to address in order to produce a good normalized set of tables:
  - Evaluate PK Assignments
  - Evaluate Naming Conventions
  - Refine Attribute Atomicity
  - Identify New Attributes
  - Identify New Relationships
  - Refine Primary Keys as Required for Data Granularity
  - Maintain Historical Accuracy
  - Evaluate Using Derived Attributes



# Improving the Design (continued)

**FIGURE 7.6** The completed database

Database name: Ch07\_ConstructCo

Table name: PROJECT

PROJ_NUM	PROJ_NAME	EMP_NUM
15	Evergreen	105
18	Amber Wave	104
22	Rolling Tide	113
25	Starflight	101

Table name: JOB

JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
500	Programmer	€28.24
501	Systems Analyst	€76.43
502	Database Designer	€82.95
503	Electrical Engineer	€66.76
504	Mechanical Engineer	€53.64
505	Civil Engineer	€44.07
506	Clerical Support	€21.23
507	DSS Analyst	€36.30
508	Applications Designer	€38.00
509	Bio Technician	€27.29
510	General Support	€14.50



## Improving the Design (continued)

Table name: ASSIGNMENT

ASSIGN_NUM	ASSIGN_DATE	PROJ_NUM	EMP_NUM	ASSIGN_HOURS	ASSIGN_CHG_HOUR	ASSIGN_CHARGE
1001	04-Mar-06	15	103	2.60	€67.55	€175.63
1002	04-Mar-06	18	118	1.40	€14.50	€20.30
1003	05-Mar-06	15	101	3.60	€82.95	€298.62
1004	05-Mar-06	22	113	2.50	€38.00	€95.00
1005	05-Mar-06	15	103	1.90	€67.55	€128.35
1006	05-Mar-06	25	115	4.20	€76.43	€321.01
1007	05-Mar-06	22	105	5.20	€82.95	€431.34
1008	05-Mar-06	25	101	1.70	€82.95	€141.02
1009	05-Mar-06	15	105	2.00	€82.95	€165.90
1010	06-Mar-06	15	102	3.80	€76.43	€290.43
1011	06-Mar-06	22	104	2.60	€76.43	€198.72
1012	06-Mar-06	15	101	2.30	€82.95	€190.79





# Improving the Design (continued)

ASSIGN_NUM	ASSIGN_DATE	PROJ_NUM	EMP_NUM	ASSIGN_HOURS	ASSIGN_CHG_HOUR	ASSIGN_CHARGE
1013	06-Mar-06	25	114	1.80	€38.00	€68.40
1014	06-Mar-06	22	111	4.00	€21.23	€84.92
1015	06-Mar-06	25	114	3.40	€38.00	€129.20
1016	06-Mar-06	18	112	1.20	€36.30	€43.56
1017	06-Mar-06	18	118	2.00	€14.50	€29.00
1018	06-Mar-06	18	104	2.60	€76.43	€196.72
1019	06-Mar-06	15	103	3.00	€67.55	€202.65
1020	07-Mar-06	22	105	2.70	€82.95	€223.97
1021	08-Mar-06	25	108	4.20	€76.43	€321.01
1022	07-Mar-06	25	114	5.80	€38.00	€220.40
1023	07-Mar-06	22	106	2.40	€28.24	€67.78

Table Name: EMPLOYEE

EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE
101	News	John	G	08-Nov-00	502
102	Senior	David	H	12-Jul-89	501
103	Arbough	June	E	01-Dec-97	503
104	Ramoras	Anne	K	15-Nov-88	501
105	Johnson	Alice	K	01-Feb-94	502
106	Smithfield	William		22-Jun-05	500
107	Alonzo	Maria	D	10-Oct-94	500
108	Washington	Ralph	B	22-Aug-89	501
109	Smith	Larry	W	18-Jul-99	501
110	Olenko	Gerald	A	11-Dec-96	505
111	Wabash	Geoff	B	04-Apr-89	506
112	Smithson	Darlene	M	23-Oct-95	507
113	Joenbrood	Delbert	K	15-Nov-94	508
114	Jones	Anneise		20-Aug-91	508
115	Bawangi	Travis	B	25-Jan-90	501
116	Pratt	Gerald	L	05-Mar-95	510
117	Williamson	Angie	H	19-Jun-94	509
118	Frommer	James	J	04-Jan-06	510

## Surrogate Key Considerations

- When primary key is considered to be unsuitable, designers use surrogate keys
- Data entries in Table 7.3 are inappropriate because they duplicate existing records
  - Yet there has been no violation of either entity integrity or referential integrity



## Surrogate Key Considerations (continued)

**TABLE 7.3** Duplicate entries in the job table

JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
511	Programmer	€26.66
512	Programmer	€26.66



## The Boyce-Codd Normal Form (BCNF)

- Every determinant in table is a candidate key
  - Has same characteristics as primary key, but for some reason, not chosen to be primary key
- When table contains only one candidate key, the 3NF and the BCNF are equivalent
- BCNF can be violated only when table contains more than one candidate key





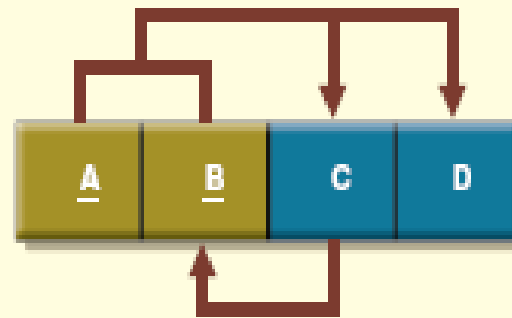
# The Boyce-Codd Normal Form (BCNF) (continued)

- Most designers consider the BCNF as special case of 3NF
- Table is in 3NF when it is in 2NF and there are no transitive dependencies
- Table can be in 3NF and fails to meet BCNF
  - No partial dependencies, nor does it contain transitive dependencies
  - A nonkey attribute is the determinant of a key attribute



# The Boyce-Codd Normal Form (BCNF) (continued)

**FIGURE 7.7** A table that is in 3NF but not in BCNF



# The Boyce-Codd Normal Form (BCNF) (continued)

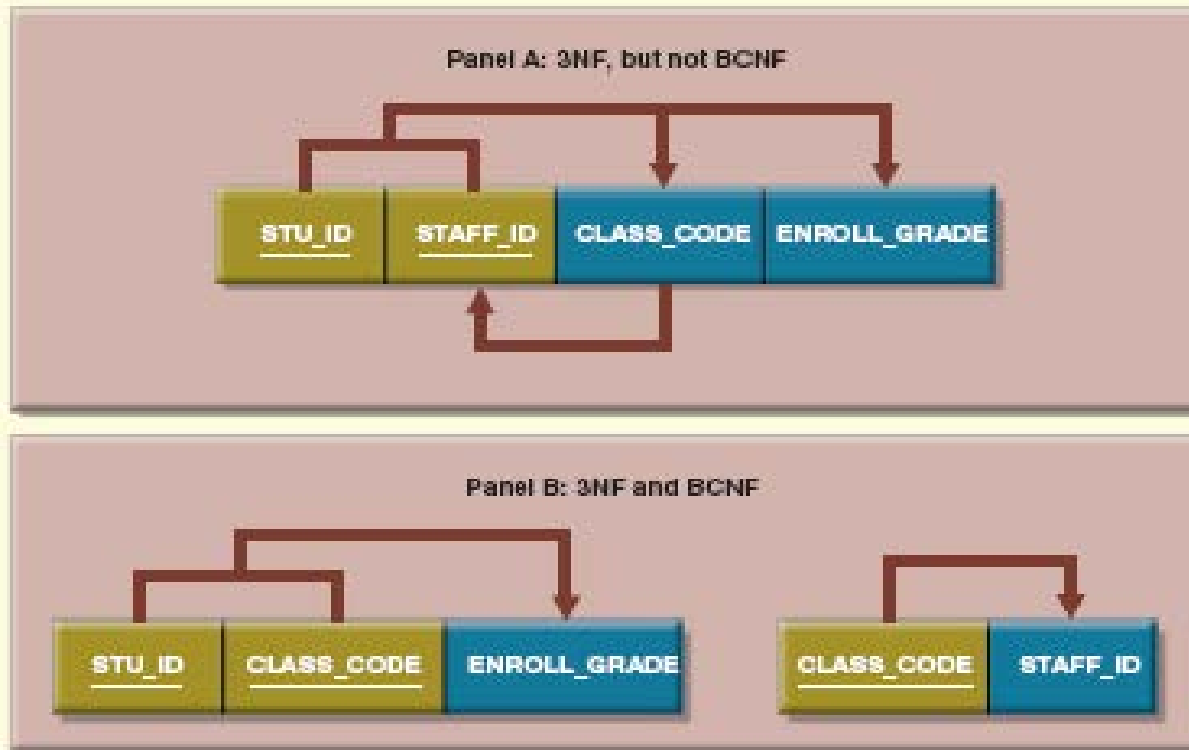
**TABLE 7.4** Sample data for a BCNF conversion—amended]

STU_ID	STAFF_ID	CLASS_CODE	ENROLL_GRADE
125	25	21334	A
125	20	32456	C
135	20	28458	B
144	25	27563	C
144	20	32456	B



# The Boyce-Codd Normal Form (BCNF) (continued)

**FIGURE 7.9** Another BCNF decomposition





## Fourth Normal Form (4NF)

- Table is in fourth normal form (4NF) when both of the following are true:
  - It is in 3NF
  - Has no multiple sets of multivalued dependencies
- 4NF is largely academic if tables conform to following two rules:
  - All attributes must be dependent on primary key, but independent of each other
  - No row contains two or more multivalued facts about an entity



# Fourth Normal Form (4NF) (continued)

**FIGURE 7.10** Tables with multivalued dependencies

Database name: Ch07\_Service

Table name: VOLUNTEER\_V1

EMP_NUM	ORG_CODE	ASSIGN_NUM
10123	RC	1
10123	UW	3
10123		4

Table name: VOLUNTEER\_V2

EMP_NUM	ORG_CODE	ASSIGN_NUM
10123	RC	
10123	UW	
10123		1
10123		3
10223		4

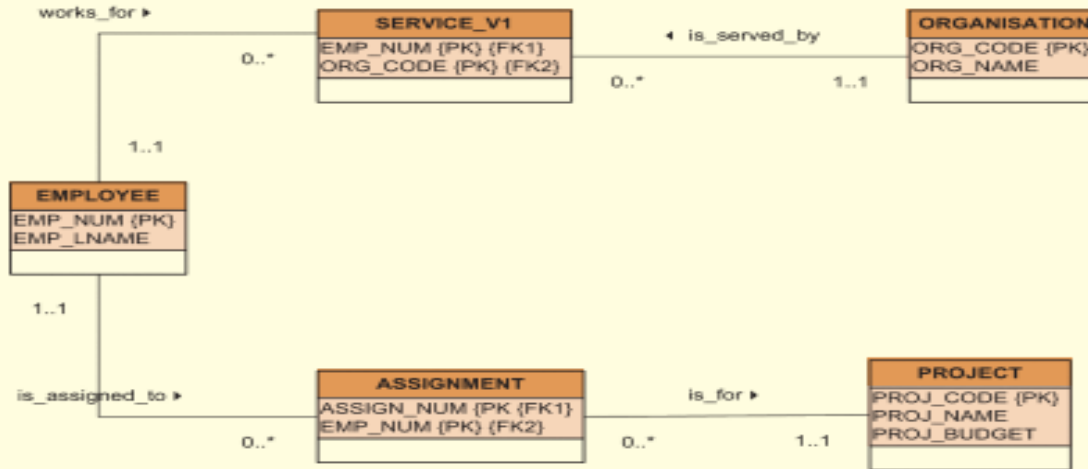
Table name: VOLUNTEER\_V3

EMP_NUM	ORG_CODE	ASSIGN_NUM
10123	RC	1
10123	RC	3
10123	UW	4

# Fourth Normal Form (4NF)

**FIGURE 7.11** A set of tables in 4NF

Relational Diagram



Database name: Ch07\_Service

Table name: EMPLOYEE

EMP_NUM	EMP_LNAME
10121	Rogers
10122	O'Leery
10123	Panera
10124	Johnson

Table name: PROJECT

PROJ_CODE	PROJ_NAME	PROJ_BUDGET
1	BeThere	€808,363.55
2	BlueMoon	€15,956,900.32
3	GreenThumb	€2,555,220.24
4	GoFast	€4,482,460.00
5	GoSlow	€791,975.00

Table name: ORGANIZATION

ORG_CODE	ORG_NAME
RC	Red Cross
UW	United Way
WF	Wildlife Fund



## Fourth Normal Form (4NF)

Table name: ASSIGNMENT

ASSIGN_NUM	EMP_NUM	PROJ_CODE
1	10123	1
2	10121	2
3	10123	3
4	10123	4
5	10121	1
6	10124	2
7	10124	3
8	10124	5

Table name: SERVICE\_V1

EMP_NUM	ORG_CODE
10123	RC
10123	UW
10123	WF



# Normalization and Database Design

- Normalization should be part of design process
- Make sure that proposed entities meet required normal form before table structures are created
- Many real-world databases have been improperly designed or burdened with anomalies if improperly modified during course of time
- You may be asked to redesign and modify existing databases



# Normalization and Database Design (continued)

- ER diagram
  - Provides big picture, or macro view, of an organization's data requirements and operations
  - Created through an iterative process
    - Identifying relevant entities, their attributes and their relationship
    - Use results to identify additional entities and attributes



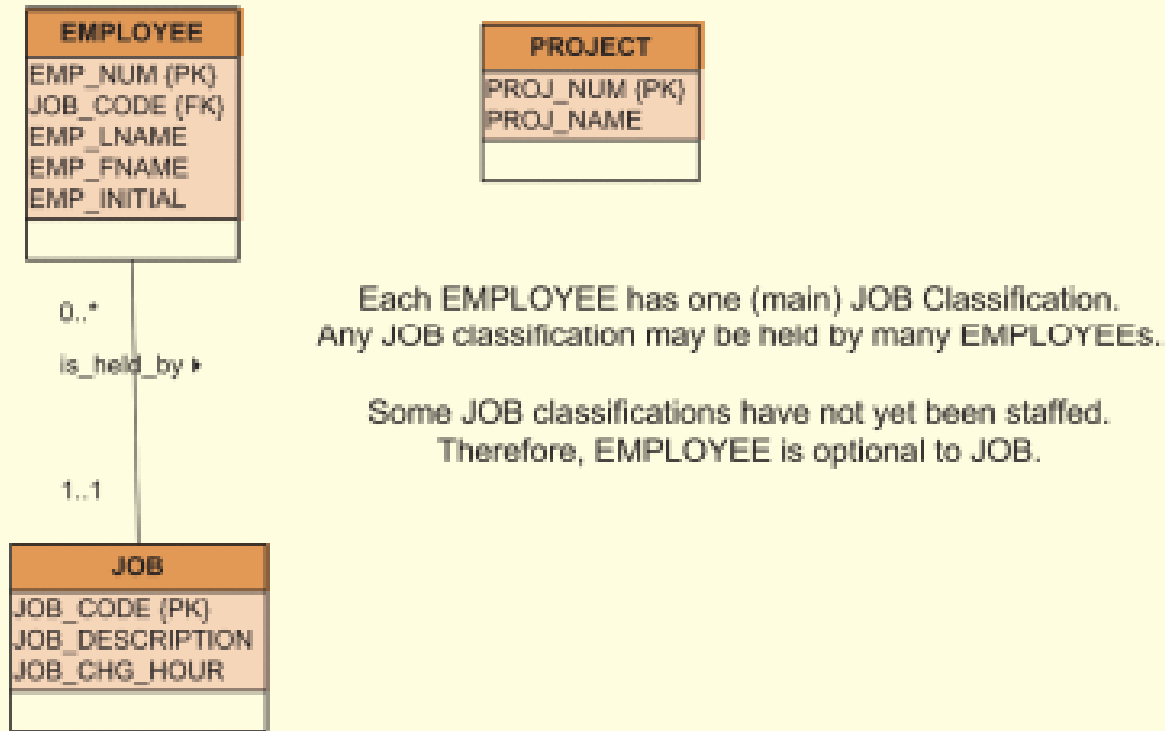
# Normalization and Database Design (continued)

- Normalization procedures
  - Focus on characteristics of specific entities
  - Represents micro view of entities within ER diagram
- Difficult to separate normalization process from ER modeling process
- Two techniques should be used concurrently



# Normalization and Database Design (continued)

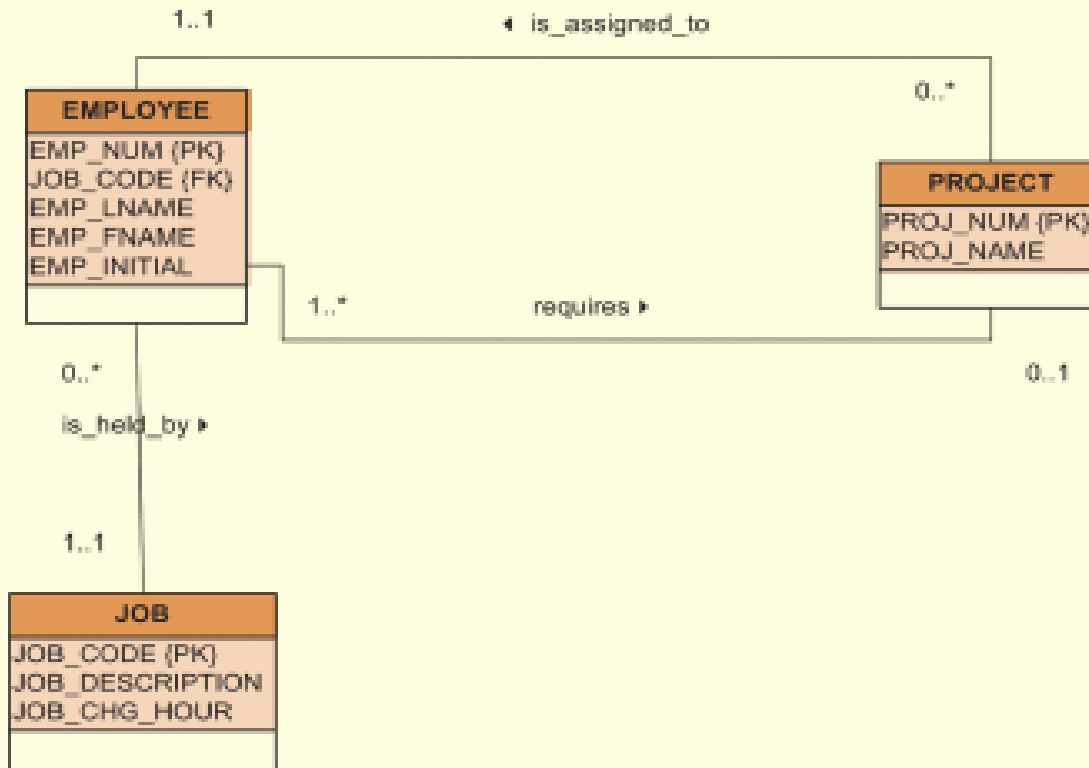
**FIGURE 7.13** Modified contracting company ERD





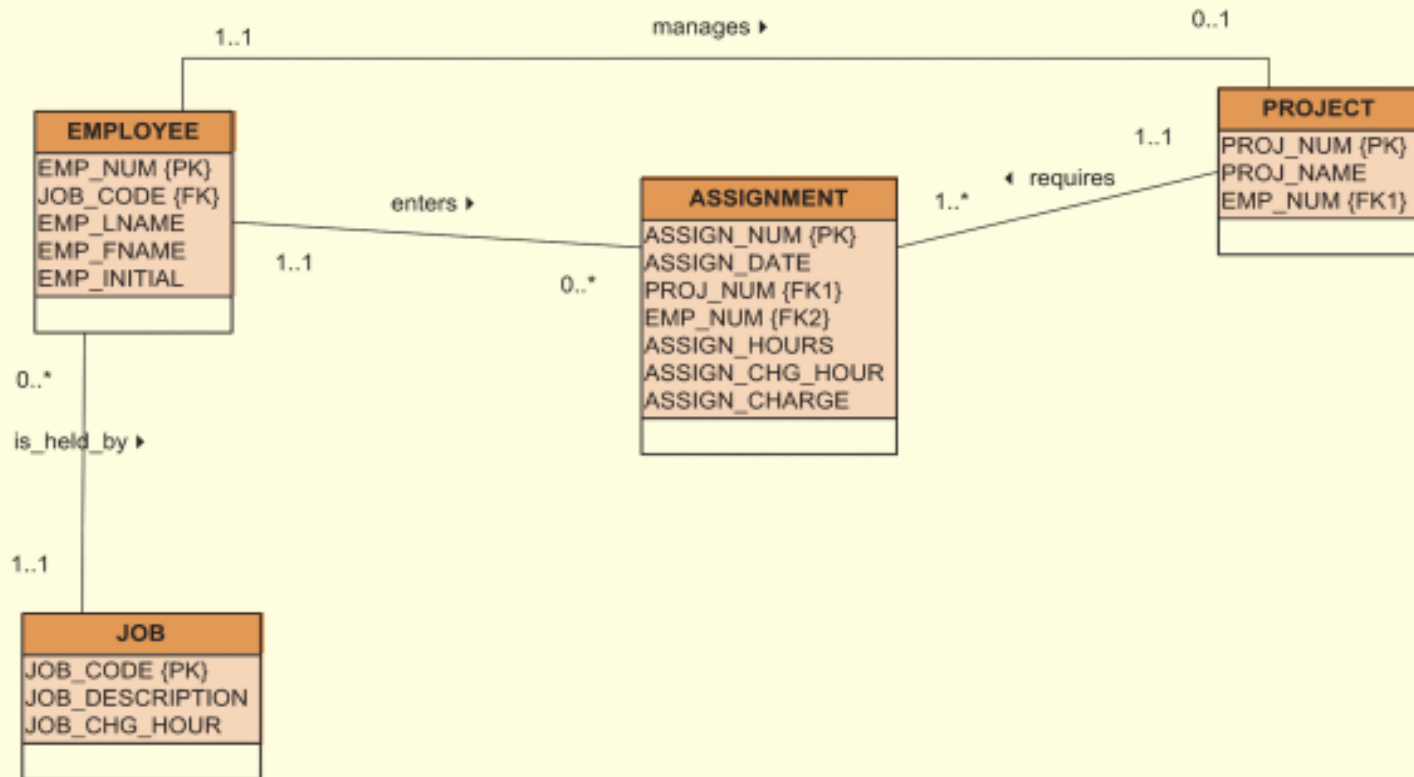
# Normalization and Database Design (continued)

**FIGURE 7.14** Modified contracting company ERD



# Normalization and Database Design (continued)

FIGURE 7.15 Final contracting company ERD



# Normalization and Database Design (continued)

**FIGURE 7.16** The implemented database

Database name: Ch07\_ConstructCo

Table name: EMPLOYEE

EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE
101	News	John	G	08-Nov-00	502
102	Senior	David	H	12-Jul-89	501
103	Arbough	June	E	01-Dec-97	503
104	Ramoras	Anne	K	15-Nov-88	501
105	Johnson	Alice	K	01-Feb-94	502
106	Smithfield	William		22-Jun-05	500
107	Alonzo	Maria	D	10-Oct-94	500
108	Washington	Ralph	B	22-Aug-89	501
109	Smith	Larry	W	18-Jul-99	501
110	Olenko	Gerald	A	11-Dec-96	505
111	Wabash	Geoff	B	04-Apr-89	506
112	Smithson	Darlene	M	23-Oct-95	507
113	Joebrood	Delbert	K	15-Nov-94	508
114	Jones	Annelise		20-Aug-91	508
115	Bawangi	Travis	B	25-Jan-90	501
116	Pratt	Gerald	L	05-Mar-95	510



# Normalization and Database Design (continued)

EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_INITIAL	EMP_HIREDATE	JOB_CODE
117	Williamson	Angie	H	19-Jun-94	509
118	Frommer	James	J	04-Jan-06	510

Table name: JOB

JOB_CODE	JOB_DESCRIPTION	JOB_CHG_HOUR
500	Programmer	€28.24
501	Systems Analyst	€76.43
502	Database Designer	€82.95
503	Electrical Engineer	€68.78
504	Mechanical Engineer	€53.64
505	Civil Engineer	€44.07
506	Clerical Support	€21.23
507	DSS Analyst	€38.30
508	Applications Designer	€38.00
509	Bio Technician	€27.29
510	General Support	€14.50

Table name: PROJECT

PROJ_NUM	PROJ_NAME	EMP_NUM
15	Evergreen	105
18	Amber Wave	104
22	Rolling Tide	113
25	Starlight	101

Table name: ASSIGNMENT

ASSIGN_NUM	ASSIGN_DATE	PROJ_NUM	EMP_NUM	ASSIGN_HOURS	ASSIGN_CHG_HOUR	ASSIGN_CHARGE
1001	04-Mar-06	15	103	2.80	€67.55	€175.83
1002	04-Mar-06	18	118	1.40	€14.50	€20.30
1003	05-Mar-06	15	101	3.80	€82.95	€298.82
1004	05-Mar-06	22	113	2.50	€38.00	€95.00
1005	05-Mar-06	15	103	1.90	€67.55	€128.35
1006	05-Mar-06	25	115	4.20	€76.43	€321.01
1007	05-Mar-06	22	105	5.20	€82.95	€431.34
1008	05-Mar-06	25	101	1.70	€82.95	€141.02
1009	05-Mar-06	15	105	2.00	€82.95	€165.90
1010	06-Mar-06	15	102	3.80	€76.43	€290.43

# Normalization and Database Design (continued)

ASSIGN_ NUM	ASSIGN_ DATE	PROJ_ NUM	EMP_ NUM	ASSIGN_ HOURS	ASSIGN_CHG_ HOUR	ASSIGN_ CHARGE
1011	06-Mar-06	22	104	2.60	€76.43	€198.72
1012	06-Mar-06	15	101	2.30	€82.95	€190.79
1013	06-Mar-06	25	114	1.80	€38.00	€68.40
1014	06-Mar-06	22	111	4.00	€21.23	€84.92
1015	06-Mar-06	25	114	3.40	€38.00	€129.20
1016	06-Mar-06	18	112	1.20	€36.30	€43.56
1017	06-Mar-06	18	118	2.00	€14.50	€29.00
1018	06-Mar-06	18	104	2.60	€76.43	€198.72
1019	06-Mar-06	15	103	3.00	€67.55	€202.65
1020	07-Mar-06	22	105	2.70	€82.95	€223.97
1021	08-Mar-06	25	108	4.20	€76.43	€321.01
1022	07-Mar-06	25	114	5.80	€38.00	€220.40
1023	07-Mar-06	22	106	2.40	€28.24	€67.78

## Denormalization

- Creation of normalized relations is important database design goal
- Processing requirements should also be a goal
- If tables decomposed to conform to normalization requirements:
  - Number of database tables expands



## Denormalization (continued)

- Joining the larger number of tables takes additional input/output (I/O) operations and processing logic, thereby reducing system speed
- Conflicts between design efficiency, information requirements, and processing speed are often resolved through compromises that may include denormalization



## Denormalization (continued)

- Unnormalized tables in production database tend to suffer from these defects:
  - Data updates are less efficient because programs that read and update tables must deal with larger tables
  - Indexing is more cumbersome
  - Unnormalized tables yield no simple strategies for creating virtual tables known as views





## Denormalization (continued)

- Use denormalization cautiously
- Understand why—under some circumstances—unnormalized tables are better choice



## Summary

- Normalization is technique used to design tables in which data redundancies are minimized
- First three normal forms (1NF, 2NF, and 3NF) are most commonly encountered
- Table is in 1NF when all key attributes are defined and when all remaining attributes are dependent on primary key



## Summary (continued)

- Table is in 2NF when it is in 1NF and contains no partial dependencies
- Table is in 3NF when it is in 2NF and contains no transitive dependencies
- Table that is not in 3NF may be split into new tables until all of the tables meet 3NF requirements
- Normalization is important part—but only part—of design process

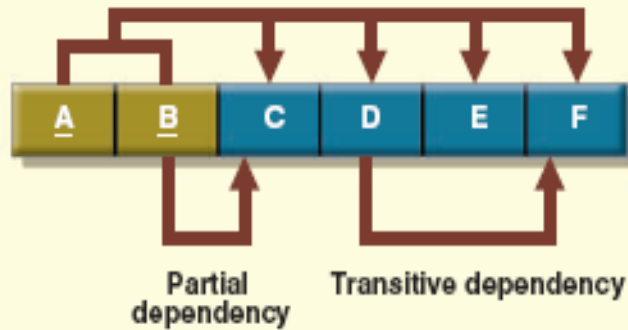




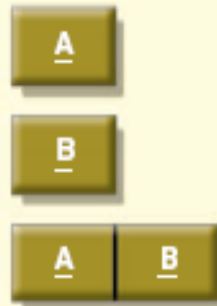
# Summary (continued)

**FIGURE 7.17**

The initial 1NF structure



Step 1: Write each PK component on a separate line; then write the original (composite) PK on the last line.





# Summary (continued)

**FIGURE 7.18** Identifying possible PK attributes

Step 2: Place all dependent attributes with the PK attributes identified in Step 1.

A No attributes are dependent on A. Therefore, A does not become a PK for a new table structure.

B C This table is in 3NF because it is in 2NF (no partial dependencies) and it contains no transitive dependencies.

A B D E F This table is in 2NF because it contains a transitive dependency.

Transitive dependency

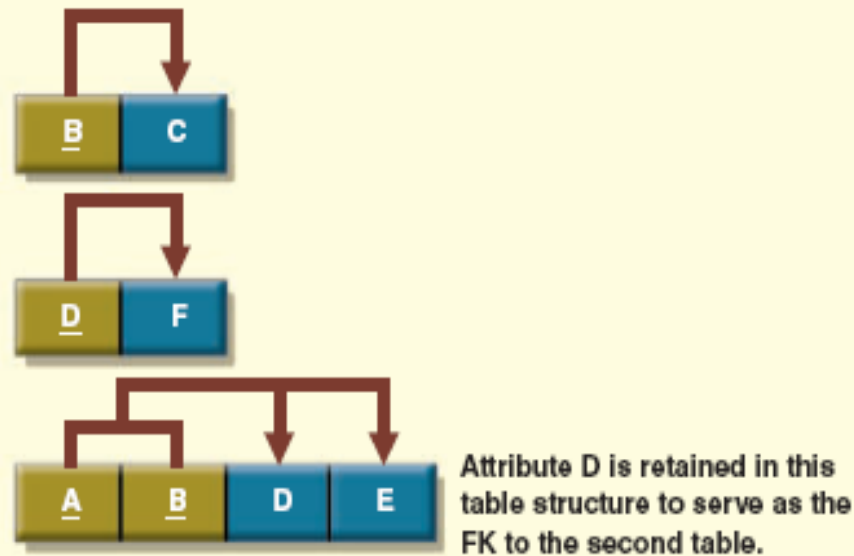
## Summary (continued)

**FIGURE 7.19**

Table structures based on the selected PKs

**Step 3: Remove all transitive dependencies identified in Step 2 and retain all 3NF structures.**

All tables are in 3NF because they are in 2NF (no partial dependencies) and they do not contain transitive dependencies.





## Summary (continued)

- Table in 3NF may contain multivalued dependencies that produce either numerous null values or redundant data
- It may be necessary to convert 3NF table to fourth normal form (4NF) by
  - Splitting table to remove multivalued dependencies
- Tables are sometimes denormalized to yield less I/O which increases processing speed

